

Age-of-Information-based Scheduling in Multiuser Uplinks with Stochastic Arrivals: A POMDP Approach

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Background

- The information freshness has become an increasingly important performance metric in this era of the Internet of Things (IoT).
- The concept of the **age of information (AoI)** has been proposed to measure the information freshness from the monitor's perspective.
- In this paper, we focus on scheduling problems of minimizing the network-wide AoI in **multiuser uplinks with stochastic arrivals**.

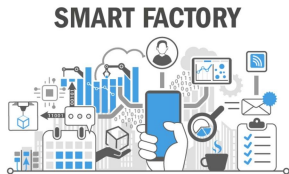


Figure: Smart Factory



Figure: Smart Healthcare



Figure: Smart Transport

Multiuser Downlinks with Stochastic Arrivals

- There is **no uncertainty about the status update arrivals from the monitor's perspective** since status updates arrive at the monitor.

What about **multiuser uplinks with stochastic arrivals**?

Weakness of Existing Work

- Most existing work assumed end nodes **used extra feedback** to report their status update arrivals to the monitor.
- Such feedback leads to **considerable overhead** and thus makes the corresponding scheduling policies **hard to implement**.

To combat this weakness, we assume that **there is no extra feedback**, which leads to a scheduling problem under **partial system information**.

System Model - I

- We consider a multiuser uplink system where K end nodes report their freshest status updates to a common monitor.
- In each time slot, a new status update arrives at end node $i \in \{1, 2, \dots, K\}$ with probability $\lambda_i \in (0, 1]$.

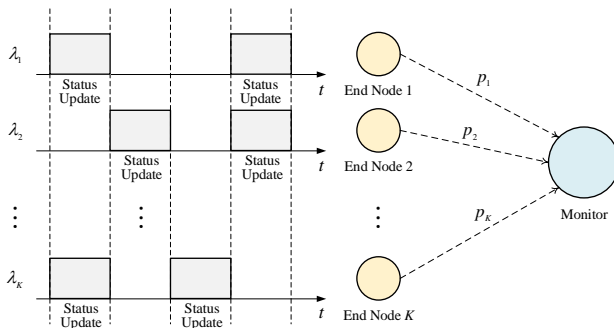


Figure: The multiuser uplink system with stochastic arrivals of status updates.

System Model - II

- At the beginning of each time slot, the monitor **schedules at most one end node** to transmit its freshest status update.
- The transmission of end node i to the monitor has a **successful probability p_i** and an **error probability $(1 - p_i)$** .

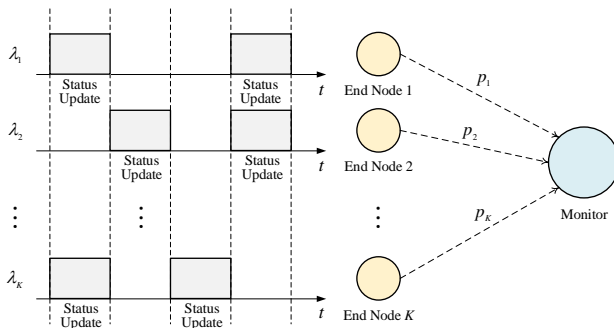
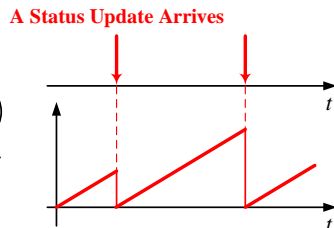
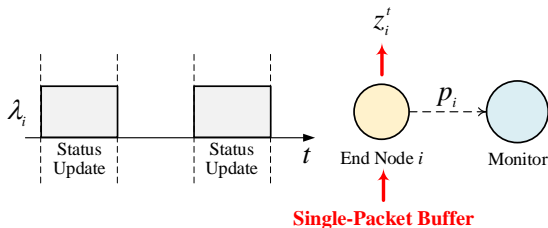


Figure: The multiuser uplink system with stochastic arrivals of status updates.

Local Age

- Define **the local age of end node i** as the time elapsed since the generation of the freshest status update at end node i .
- The evolution of **the local age of end node i** , denoted by z_i^t , is

$$z_i^{t+1} = \begin{cases} z_i^t + 1, & \text{no status update arrives in slot } t, \\ 1, & \text{a new status update arrives in slot } t. \end{cases} \quad (1)$$

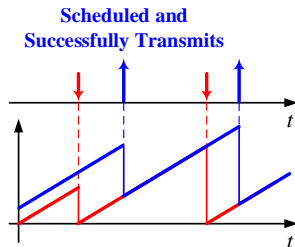
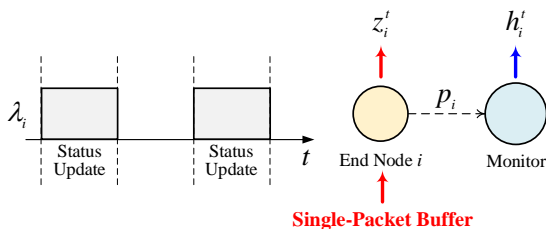


Age of Information

Age of Information

- Define **the age of information (Aol)** of end node i as the time elapsed since the generation (at end node i) of the latest received status update (at the monitor).
- The evolution of **the Aol of end node i** , denoted by h_i^t , is

$$h_i^{t+1} = \begin{cases} z_i^t + 1, & \text{scheduled and successfully transmits in slot } t \\ h_i^t + 1, & \text{otherwise.} \end{cases} \quad (2)$$



Main Contributions

Weakness of Existing Work

- Most existing work assumed end nodes **used extra feedback** to report their status update arrivals to the monitor.
- Such feedback leads to **considerable overhead** and thus makes the corresponding scheduling policies **hard to implement**.

To combat this weakness, we assume that **there is no extra feedback**, which leads to a scheduling problem under **partial system information**.

Main Contributions

- Formulate **a partially observable Markov decision process (POMDP)** to characterize the dynamic behavior of such system.
- Use a dynamic programming (DP) algorithm to attain **the optimal policy** and devise **a myopic policy** with low computation complexity.

A POMDP models an agent decision process in which it is assumed that

- the system dynamics are determined by an MDP, and
- the agent cannot directly observe the underlying state.

Definition

A POMDP can be described as a tuple $\langle \mathcal{S}, \mathcal{A}, \mathcal{O}, T, R, Z \rangle$.

- \mathcal{S} is a finite set of states.
- \mathcal{A} is a finite set of actions.
- \mathcal{O} is a finite set of observations.
- T is the state transition function, $T : \mathcal{S} \times \mathcal{A} \rightarrow \Pi(\mathcal{S})$.
- R is the reward function, $R : \mathcal{S} \times \mathcal{A} \rightarrow \mathbb{R}$.
- Z is the observation function, $Z : \mathcal{S} \times \mathcal{A} \rightarrow \Pi(\mathcal{O})$.

States

- Denote **a state of end node i** in slot t by $\mathbf{s}_i^t \triangleq [h_i^t, z_i^t]$.
 - $h_i^t \in \mathcal{T} \triangleq \{1, 2, 3, \dots\}$ is its instantaneous Aol at the monitor
 - $z_i^t \in \mathcal{T}$ is its local age
- Denote **a state of the POMDP** in slot t by $\mathbf{s}^t \triangleq [\mathbf{h}^t, \mathbf{z}^t]$.
 - $\mathbf{h}^t \triangleq [h_1^t, \dots, h_K^t] \in \mathcal{H} = \mathcal{T}^K$ is the Aol of all end nodes
 - $\mathbf{z}^t \triangleq [z_1^t, \dots, z_K^t] \in \mathcal{Z} = \mathcal{T}^K$ is the local age of all end nodes

Actions

- Denote **an action of end node i** in slot t by $a_i^t \in \{0, 1\}$.
 - If end node i is scheduled, $a_i^t = 1$; otherwise, $a_i^t = 0$.
- Denote **an action of the POMDP** in slot t by $\mathbf{a}^t \triangleq [a_1^t, \dots, a_K^t]$.
 - In the single-antenna system considered, we have $\sum_{i=1}^K a_i^t \leq 1$.

Observations

- Denote **an observation of end node i** in slot t by $\mathbf{o}_i^t \triangleq [h_i^t, \hat{z}_i^t]$.
 - h_i^t is its fully observed AoI at the monitor
 - $\hat{z}_i^t \in \{\mathcal{T}, X\}$ is its partially observed local age, where X means no observation of the local age of an end node
- Denote **an observation of the POMDP** in slot t by $\mathbf{o}^t \triangleq [\mathbf{o}_1^t, \dots, \mathbf{o}_K^t]$.

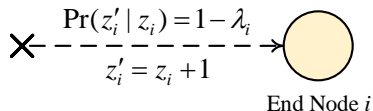
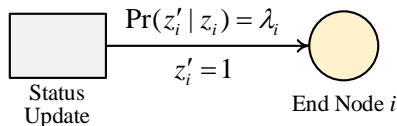
Belief State

- Denote **a belief state of the POMDP** in slot t by $\mathbf{l}^t \triangleq [\mathbf{h}^t, \mathbf{b}^t]$.
 - \mathbf{b}^t is a probability distribution over \mathcal{Z} , where $b(\mathbf{z})$ is the probability assigned to $\mathbf{z}^t = \mathbf{z}$ by $\mathbf{b}^t = \mathbf{b}$, satisfying $\sum_{\mathbf{z} \in \mathcal{Z}} b(\mathbf{z}) = 1$.
 - Note that \mathbf{h}^t is fully observable, i.e., **its belief update is always deterministic given \mathbf{h}^{t-1} , \mathbf{a}^{t-1} and \mathbf{o}^{t-1} .**

POMDP Formulation - III

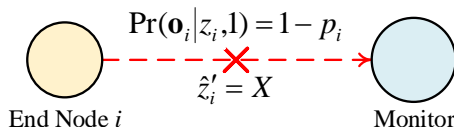
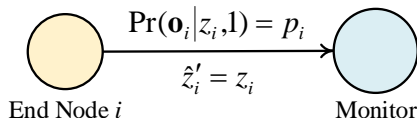
State Transition Function

- Let $\Pr(\mathbf{z}'|\mathbf{z}) = \prod_{i=1}^K \Pr(z'_i|z_i)$ denote the state transition function.



Observation Function

- Let $\Pr(\mathbf{o}|\mathbf{z}, \mathbf{a}) = \prod_{i=1}^K \Pr(\mathbf{o}_i|z_i, a_i)$ denote the observation function.



Belief Update

- When given $\mathbf{l}^t = \mathbf{l}$, $\mathbf{a}^t = \mathbf{a}$ and $\mathbf{o}^t = \mathbf{o}$, for $\forall i$, h_i^{t+1} can be updated as follows:

$$h_i^{t+1} = \begin{cases} \hat{z}_i + 1, & \text{if } \hat{z}_i \neq X, \\ h_i + 1, & \text{if } \hat{z}_i = X. \end{cases} \quad (3)$$

- When given the same condition, for $\forall \mathbf{z}' \in \mathcal{Z}$, $b^{t+1}(\mathbf{z}')$ can be updated via the Bayes' theorem:

$$b^{t+1}(\mathbf{z}') = \left\{ \sum_{\mathbf{z} \in \mathcal{Z}} \Pr(\mathbf{z}'|\mathbf{z}) \Pr(\mathbf{o}|\mathbf{z}, \mathbf{a}) b(\mathbf{z}) \right\} / \Pr(\mathbf{o}|\mathbf{l}, \mathbf{a}), \quad (4)$$

where $\Pr(\mathbf{o}|\mathbf{l}, \mathbf{a}) = \sum_{\mathbf{z} \in \mathcal{Z}} \Pr(\mathbf{o}|\mathbf{z}, \mathbf{a}) b(\mathbf{z})$.

- We denote the process by the update function $\mathbf{l}^{t+1} = f(\mathbf{l}, \mathbf{a}, \mathbf{o})$.

POMDP Formulation - V

Reward

- Define **the reward** at belief state $\mathbf{I}^t = \mathbf{I}$ as the weighted sum of the instantaneous Aol of all end nodes, i.e., $R(\mathbf{I}) \triangleq \sum_{i=1}^K \omega_i h_i$.

Policy

- Define **a decision rule** as a mapping from the belief space \mathcal{I} into the action space \mathcal{A} , i.e., $d_t : \mathcal{I} \rightarrow \mathcal{A}$.
- Define **a policy** as a sequence of decision rules, i.e., $\pi = \{d_1, \dots, d_T\}$.

Let $V^\pi(\mathbf{I})$ denote **the expected total reward** from slot 1 to slot T when $\mathbf{I}^1 = \mathbf{I}$ and the policy π is used, which can be defined by

$$V^\pi(\mathbf{I}) \triangleq \mathbb{E} \left[\sum_{t=1}^T R(\mathbf{I}^t) \mid \mathbf{I}^1 = \mathbf{I}, \pi \right] \rightarrow \text{(P): } \pi^* = \arg \min_{\pi} \frac{1}{TK} V^\pi(\mathbf{I}).$$

The Optimal Policy

The Total Expected Reward

- Let $U_t^\pi(\mathbf{I})$ denote the total expected reward from slot t to slot T when $\mathbf{I}^t = \mathbf{I}$ and the policy π is used, which can be defined by

$$U_t^\pi(\mathbf{I}) \triangleq \mathbb{E}\left[\sum_{\tau=t}^T R(\mathbf{I}^\tau) | \mathbf{I}^t = \mathbf{I}, \pi\right] \rightarrow (\text{P}): \pi^* = \arg \min_{\pi} \frac{1}{TK} U_1^\pi(\mathbf{I}).$$

Dynamic Programming

- For problem (P), we have the following bellman equation:

$$U_t^*(\mathbf{I}) = R(\mathbf{I}) + \max_{\mathbf{a} \in \mathcal{A}} \sum_{\mathbf{o} \in \mathcal{O}} \Pr(\mathbf{o} | \mathbf{I}, \mathbf{a}) U_{t+1}^*(f(\mathbf{I}, \mathbf{a}, \mathbf{o})), \quad (5)$$

for each $t \in \{1, \dots, T-1\}$ and $\mathbf{I} \in \mathcal{I}_t$.

- The minimal EWSAol given $\mathbf{I}^1 = \mathbf{I}$ can be computed by $U_1^*(\mathbf{I})/TK$.

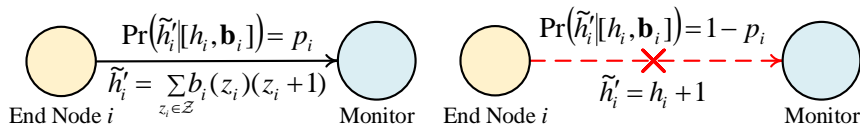
A Myopic Policy - I

Belief State (Sufficient Statistics)

- The monitor can only maintain **probability distributions of the local age of each end node**, denoted by $\mathbb{B}^t = [\mathbf{b}_1^t, \mathbf{b}_2^t, \dots, \mathbf{b}_K^t]$.
- The belief state of the POMDP can be redefined as $\mathbb{I}^t \triangleq [\mathbf{h}^t, \mathbb{B}^t]$.

The One-Step Expected Aol

- Denote the expected Aol in the next slot of end node i by $\mathbb{E}(h'_i)$.
 - $\mathbb{E}(h'_i) = (1 - a_i)(h_i + 1) + a_i^t \left(p_i \sum_{z_i \in \mathcal{T}} b_i(z_i)(z_i + 1) + (1 - p_i)(h_i + 1) \right)$



The One-Step Expected Reward

- Given belief state $\mathbb{I}^t = \mathbb{I}$, if action $\mathbf{a}^t = \mathbf{a}$ is chosen in slot t , the **one-step expected reward** of the POMDP is given by

$$\hat{R}(\mathbb{I}, \mathbf{a}) = \sum_{i=1}^K \omega_i \left[(1 - a_i)(h_i + 1) + a_i^t \left(p_i \sum_{z_i \in \mathcal{T}} b_i(z_i)(z_i + 1) + (1 - p_i)(h_i + 1) \right) \right]$$

The Myopic Policy

- Given belief state $\mathbb{I}^t = \mathbb{I}$, the myopic policy can be given by

$$d_t(\mathbb{I}^t) = \arg \min_{\mathbf{a} \in \mathcal{A}} \hat{R}(\mathbb{I}, \mathbf{a}). \quad (6)$$

Successful Transmission Probability

- The successful transmission probability of end node i is computed by

$$p_i = 1 - \Pr(r_i < r_{th}) \stackrel{(a)}{=} \exp\left(-\frac{d_i^T(2^{r_{th}} - 1)}{\text{SNR}}\right). \quad (7)$$

Benchmarking Policies

- Myopic policy (full knowledge), MPF:** Given $\mathbf{s}^t = \mathbf{s}$, the monitor schedules end node i^* satisfying $i^* = \arg \max_{\forall i} \omega_i p_i (h_i - z_i)$.
- MaxAol policy, MAP:** Given $\mathbf{h}^t = \mathbf{h}$, the monitor schedules end node i^* satisfying $i^* = \arg \max_{\forall i} \omega_i p_i h_i$.
- Randomized policy, RDP:** The monitor schedules end node i with probability $\mu_i = \sqrt{\omega_i / p_i} / (\sum_{j=1}^K \sqrt{\omega_j / p_j})$, $\forall i$.

Simulation - II

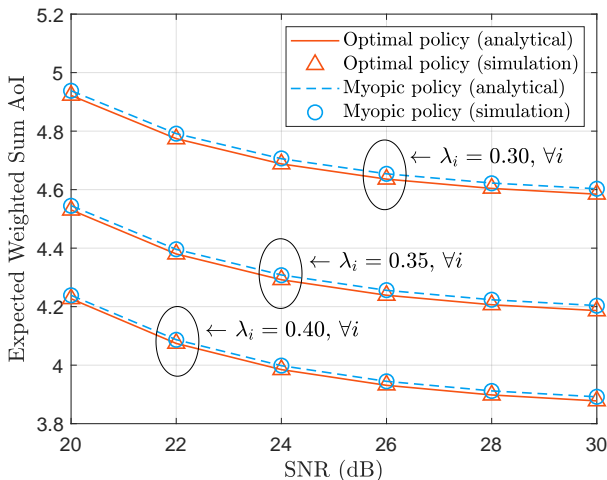


Figure: The EWSAoI as a function of **the SNR** for different configurations when $K = 2$, $T = 50$, $D = 15$, and $\omega_i = 1, \forall i$.

Simulation - III

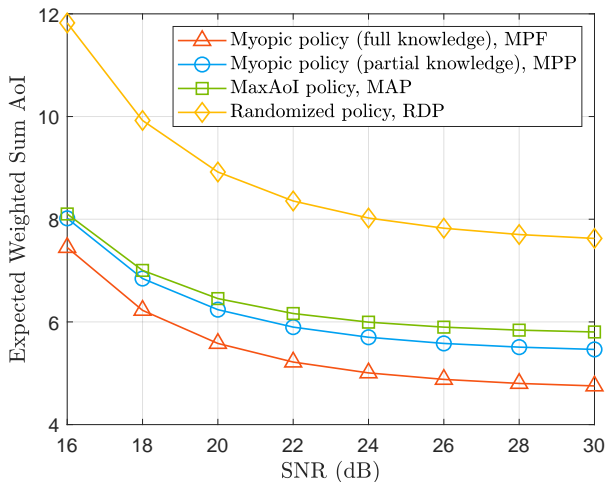


Figure: The EWSAoI as a function of the SNR when $K = 5$, $T = 10^6$, $D = 50$, $\lambda_i = 0.4$ and $\omega_i = 1, \forall i$.

Simulation - IV

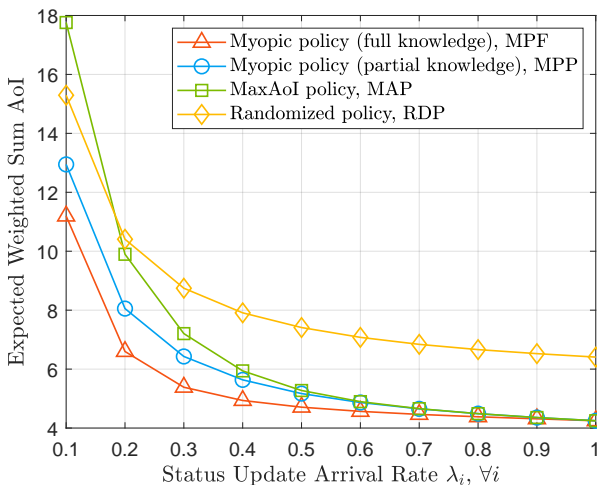


Figure: The EWSAoI as a function of **the status update arrival rate** when $K = 5$, $T = 10^6$, $D = 50$, $\text{SNR} = 25 \text{ dB}$ and $\omega_i = 1, \forall i$.

Conclusion

System Model

In this paper, we have investigated **age-of-information-based scheduling in multiuser uplinks with stochastic arrivals**.

POMDP Formulation and Solutions

- Formulate a **partially observable Markov decision process (POMDP)** to characterize the dynamic behavior of such system.
- Use a dynamic programming (DP) algorithm to attain **the optimal policy** and devise a **myopic policy** with low computation complexity.

Simulation

The performance of the myopic policy approaches that of the optimal policy and is superior to that of the benchmarking policy.